Sub-Identity: Security for Mere Mortals Phil Snowberger and Douglas Thain

Motivation

Users aren't generally capable, or inclined, to ensure the correctness of software that they download and run on their own machines. Users cannot assume any level of security on their desktops unless either they or their administrator takes some measure to compartmentalize, or contain, potentially rogue code.

Concept

• People have an intuitive understanding of familial and corporate hierarchies, in that they have a sense that higher-ups have control over those lower down.

• We propose to introduce *sub-identities* into the operating system. Each new sub-identity has a meaningful name and can be used to enforce access control, perform auditing, or simply isolate sub-processes from each other.

Current Mandatory Access Control (MAC) solutions such as SELinux are quite technically involved, requiring extensive hand-tuning by security experts. Other solutions such as traditional sandboxing and virtual machines require a high level of user sophistication to operate correctly. As effective as they promise to be, they lack intuitiveness and simplicity.

The Ideal Model, A Lofty Goal

Ideally, every user would be "effectively root" to all of its sub-users. Users should be able to:

- send signals to processes owned by their sub-users.
- implicitly access files belonging to their sub-users.
- debug processes belonging to their sub-users.

Unix doesn't happen to allow these operations, but some support for sub-identities can be "dropped in" to an existing Linux system.

A Toolkit Approximation

We have implemented an approximation of the sub-identity model in the form of a toolkit of setuid binaries that can be installed on any Linux system.

/etc/subusers root:alice,psnowber alice:betty,browser,email browser:webapp email:exec-attach



Relative Complexity

	Sub-Identity (toolkit)	Virtual Machine (VMWare)	Sandbox (Janus)
Lines of Code	885	Millions?	8743
Ops to set up	Install toolkit Call <i>subuseradd</i>	Create new image Install new OS into image Set up network bridge Install program	Install kernel module Identify system calls Define minimal policy
Ops to run	Call subusersudo	Run program in image	Run program in sandbox Iteratively update policy, if insufficient
Ops to retrieve	Call subuserchown	Extract through network	none

Here we show the identity relationships in a hypothetical system. Alice has a normal system account, whereas Charles and David authenticate through Kerberos. Each user has a normal entry in the /etc/passwd file.

The relative complexity of various compartmentalization solutions is shown. Steps in bold text require root privileges.

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A "Safe Window" is shown running an

Applications

Index of /~psnowber/bin - Mozilla Firefo	x	parrot identity box command
Index of /~psnowber/bin - Mozilla Firefo File Edit View Go Bookmark Index of /~psnowber/bin Here is a copy of /bin/bash: Den Link in New Window Open Link in New Tab Bookmark This Link Save Link As Send Link		<pre>parrot_identity_box_command 18:33:33 http://www.nd.edu/%7Epsnowber/bin/bash => 'bash' Resolving www.nd.edu 129.74.250.90 Connecting to www.nd.edu[129.74.250.90]:80 connected. HTTP request sent, awaiting response 200 OK Length: 616,312 [text/plain] 100%[=======>] 616,312K/s 18:33:33 (4.90 MB/s) - 'bash' saved [616,312/616,312] bash-3.00\$ whoami www.nd.edu bash-3.00\$ whoami www.nd.edu bash-3.00\$ pwd /tmp/home.106563/www.nd.edu bash-3.00\$ pwd /tmp/home.106563/www.nd.edu bash-3.00\$ pwd /tmp/home.106563/www.nd.edu bash-3.00\$ pwd /tmp/home.106563/www.nd.edu bash-3.00\$ pwd /tmp/home.106563/www.nd.edu bash-3.00\$ pwd /afs/nd.edu/user35/psnowber bash-3.00\$ pwd /afs/nd.edu/user35/psnowber ba</pre>
Copy Link Location	wget	
Properties	Run in an identity box	rm: cannot remove `Sent': Permission denied bash-3,00\$ ■
Done		

A quick modification to a Firefox extension allows



Xnest. Everything contained in the window is run as a sub-user. Potentially unstable or rogue X applications can be run inside it without compromising the security of the main X server. downloadable applications to be run seamlessly inside a security domain that is created on the fly. The callouts show the application being downloaded, the application under typical usage, and an unsuccessful attack being carried out. • Process hierarchy: alice's `ps` command would only show the processes inside her box.

• Disk / chroot hierarchy?



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Sub-Identities: Towards Operating System Support for Distributed System Security, Philip Snowberger and Douglas Thain, Technical Report 2005-18, University of Notre Dame, Computer Science and Engineering Dept, October 2005.

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